

Introduction to Information Science and Technology

(SI100B Python Part)

Fu Song
School of Information Science and Technology
ShanghaiTech University

Objective

To understand basic techniques for *analyzing algorithms*.

- CS140 Algorithm
- To understand the algorithms for *linear* and *binary search*.
- To understand *sorting* and know the algorithms for *selection* sort and *merge* sort.
- To understand some problems are *intractable* and others are *unsolvable*.

Roadmap

- Searching problem
 - Linear and binary search
- Recursion and Divide-and-Conquer
- Sorting problem
 - Selection sort and merge sort
- Intractability and unsolvability

Search problem

- Searching is a basic, but very important problem.
- Specification :

```
def search(x, nums):  
    # nums is a list of numbers and x is a number  
    # Returns the index in the list if x occurs  
    # or -1 if x is not in the list.
```

Built-in search methods

- Using `in`:

```
if x in nums:  
    # do something
```

- The `index` method (throws an exception if `x` is not in the list.):

```
>>> nums = [3, 1, 4, 2, 5]  
>>> nums.index(4)  
2
```

- You can implement search by using `index`, however, we want to know how `index` is implemented.

Linear search

- Pretend you were given a page full of randomly ordered numbers and were asked where 13 was.
- How would you do it?
- Scanning downward from the top to the bottom!

Linear search

```
def search(x, nums):  
    for i in range(len(nums)):  
        if nums[i] == x:  
            return i  
    return -1
```

- The Python `in` and `index` operations both implement linear searching algorithms.

Binary search

- If numbers are in ascending order, what would you do?
- Like finding some page number in a book.
- You can “throw away” some parts of the list after one comparison.

Binary search

```
def search(x, nums):  
    low = 0  
    high = len(nums) - 1  
    while low <= high:  
        mid = (low + high) / 2  
        item = nums[mid]  
        if x == item:  
            return mid  
        elif x < item:  
            high = mid - 1  
        else:  
            low = mid + 1  
    return -1
```

Comparing algorithms

- Intuitively, we might expect the linear search to work better for small lists, and binary search for longer lists. But how can we be sure?
- Method#1: Test them!
 1. $n \leq 10$, linear search faster
 2. $10 < n \leq 1000$, little difference
 3. $n > 1000$, binary search faster
 4. $n = 100,000,000$, linear 2.5s, binary .0003s

Comparing algorithms

- Method#1 is dependent on hardware (cpu, etc.) and software (compilers, etc.).
- Method#2: counting the steps!
- *How many steps are needed to find a value in a list of size n ?*
- Linear: the steps required are linearly related to n
- Binary: the steps required are related to $\log n$

Comparing algorithms

- Binary search is much faster. But why the built-in functions use linear search?
- Before carrying out binary search, you need to sort the list!
- How efficient can we solve sorting problem?
- Answer is at least $n \log n$ steps for lists sized n .

Roadmap

- Searching problem
 - Linear and binary search
- **Recursion and Divide-and-Conquer**
- Sorting problem
 - Selection sort and merge sort
- Intractability and unsolvability

Recursive version of Binary search

```
mid = (low + high) / 2
```

```
if low > high
```

```
    x is not in nums
```

```
elif x < nums[mid]
```

```
    perform binary search in nums[low]...nums[mid-1]
```

```
else
```

```
    perform binary search in nums[mid+1]...nums[high]
```

- A description of something that refers to itself is called a *recursive* definition.
- Binary search is a kind of Divide-and-conquer algorithm.

Divide-and-Conquer

- *Divide* Step: the input is partitioned into p parts.
- *Conquer* Step: q recursive call(s) of the smaller parts.
- *Combine* Step: combine the solved parts to obtain the desired output.
- In binary search:
- Divide into two parts sized $n/2$, one recursive call of the smaller part.

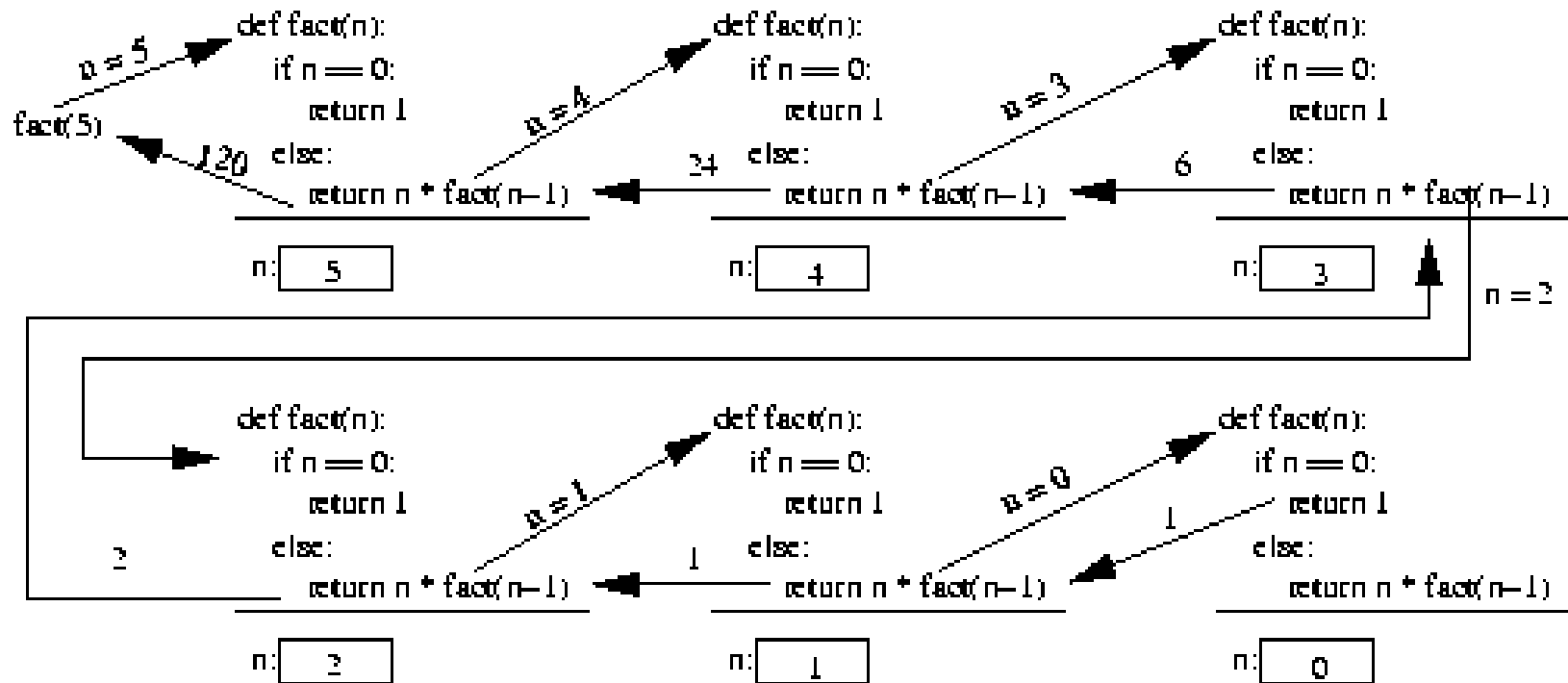
Recursion

- In mathematics, recursion is frequently used:
- $n! = n * (n-1)!$, $0! = 1$
- $\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$, $\text{fib}(1) = 1$, $\text{fib}(2) = 1$
- Two key characteristics:
 - There are one or more base cases for which no recursion is applied.
 - All chains of recursion eventually end up at one of the base cases.

Factorial function

```
def fact(n):  
    if n == 0:  
        return 1  
  
    else:  
        return n * fact(n-1)
```

Factorial calling stack



Examples

- String reversal
 - `reverse("Hello") -> "olleH"`
- Anagrams of a string
 - `anagram("abc") -> ["abc", "acb", ..., "cba"]`
- Fast exponentiation

Binary search

- ```
def recBinSearch(x, nums, low, high):
 if low > high:
 return -1
 mid = (low + high) / 2
 item = nums[mid]
 if item == x:
 return mid
 elif x < item: # Look in lower half
 return recBinSearch(x, nums, low, mid-1)
 else: # Look in upper half
 return recBinSearch(x, nums, mid+1, high)
```
- ```
def search(x, nums):  
    return recBinSearch(x, nums, 0, len(nums)-1)
```

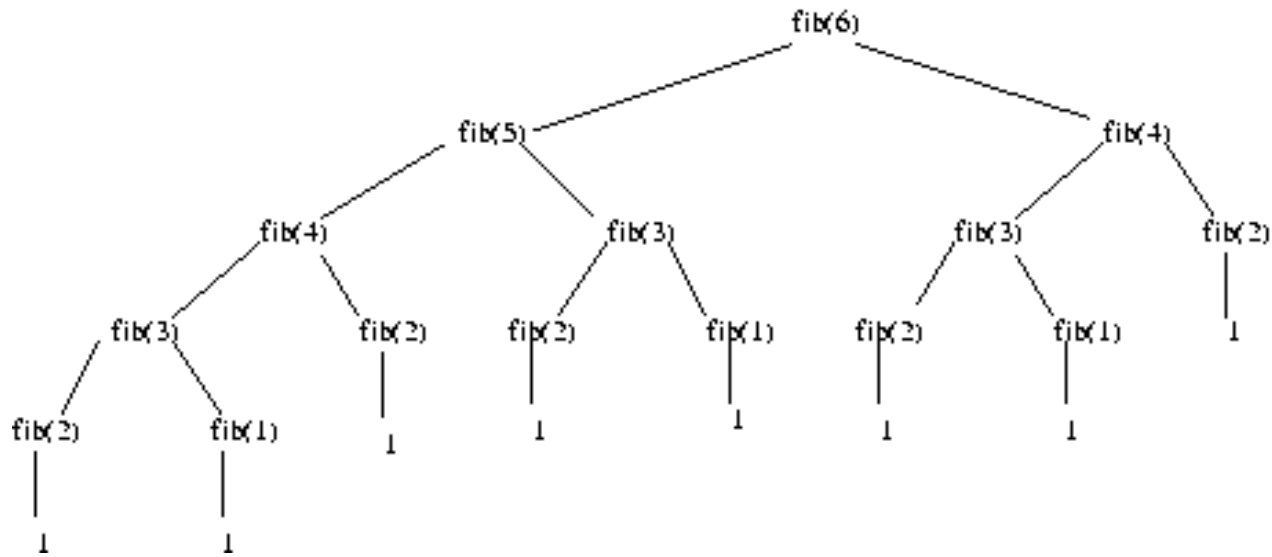
Recursion v.s. Iteration

- Anything that can be done with a loop can be done with a simple recursive function!
- Some problems that are simple to solve with recursion are quite difficult to solve with loops.
- Recursion is not always better than iterations. E.g., Fibonacci function:

$$fib(n) = \begin{cases} 1 & \text{if } n < 3 \\ fib(n-1) + fib(n-2) & \text{otherwise} \end{cases}$$

Recursion v.s. Iteration

- For Fibonacci number, the recursion version is very inefficient:



Lessons

- Recursion is another tool in your problem-solving toolbox.
- Sometimes recursion provides more elegant solution.
- At other times, when both algorithms are quite similar, use the looping solution on the basis of speed.
- Avoid the recursive solution if it is terribly inefficient, unless you can't come up with an iterative solution.

Roadmap

- Searching problem
 - Linear and binary search
- Recursion and Divide-and-Conquer
- **Sorting problem**
 - Selection sort and merge sort
- Intractability and unsolvability

Built-in sort functions

- `sorted(lst[, key = func, reverse = True])`
 - `Key` can be anonymous functions, e.g.,
`key = lambda a: a.getName()`
- `lst.sort([key= func, reverse = True])`

Naïve sorting

- Selection sort
- Loop n-1 times:
 - Find the smallest element, wap it to the correct position

```
def selSort(nums):  
    n = len(nums)  
    for bottom in range(n-1):  
        mp = bottom  
        for i in range(bottom+1, n):  
            if nums[i] < nums[mp]:  
                mp = i  
        nums[bottom], nums[mp] = nums[mp], nums[bottom]
```

Divide-and-conquer

- Merge sort
 - Divide array into two halves.
 - *Recursively* sort each half.
 - Merge two halves.
- Assume we need to merge two sorted arrays, with M , N elements respectively. Then
 - How many comparisons in *best* case?
 - How many comparisons in *worst* case?

Comparing sorts

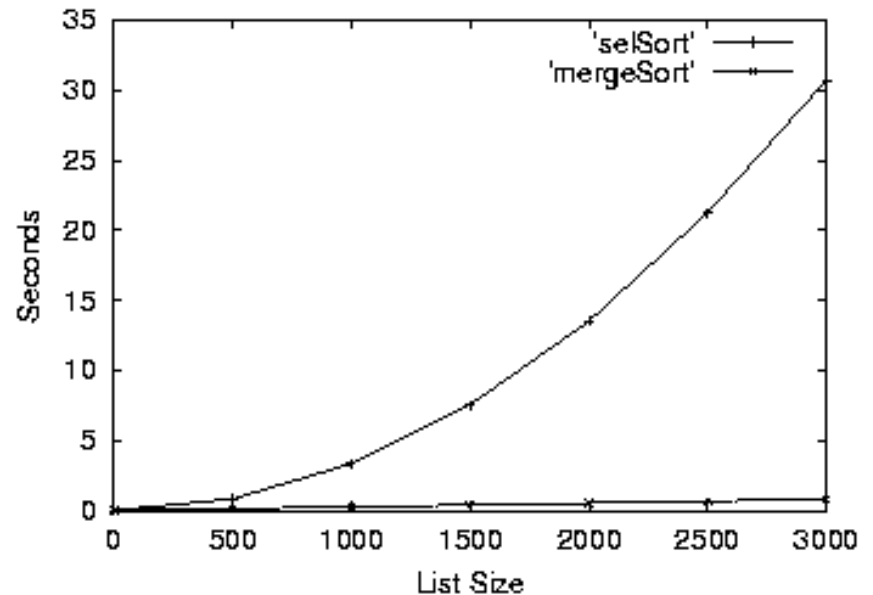
- <http://www.sorting-algorithms.com/>

- Selection sort:

- $n + (n-1) + (n-2) + (n-3) + \dots + 1$

- Merge sort:

- $n \log n$



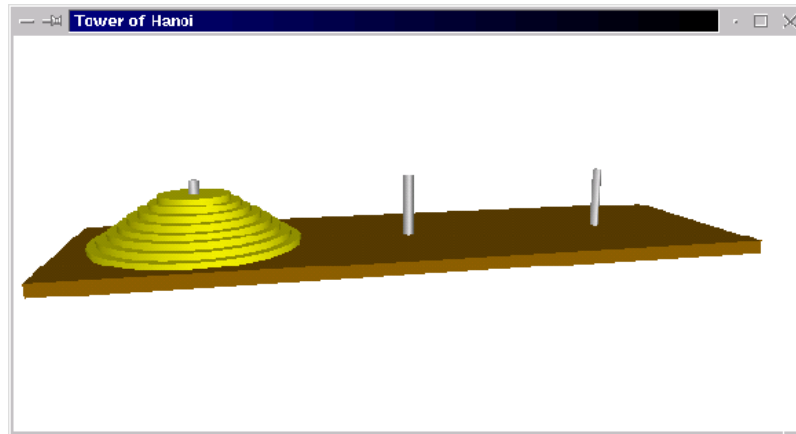
Hard Problems

- We have efficient algorithms for searching and sorting problems.
- Divide and conquer and recursion are very powerful techniques for algorithm design.
- **Not all problems** have efficient solutions!

Roadmap

- Searching problem
 - Linear and binary search
- Recursion and Divide-and-Conquer
- Sorting problem
 - Selection sort and merge sort
- **Intractability and unsolvability**

Towers of Hanoi



- ```
def moveTower(n, source, dest, temp):
 if n == 1:
 print "Move disk from", source, "to", dest+". "
 else:
 moveTower(n-1, source, temp, dest)
 moveTower(1, source, dest, temp)
 moveTower(n-1, temp, dest, source)
```
- ```
def hanoi(n):  
    moveTower(n, "A", "C", "B")
```

Analysis of Hanoi

- To solve a puzzle of size n will require 2^{n-1} steps -- an exponential time algorithm.
- For 64 disks, moving one a second, round the clock, would require *580 billion years* to complete (The current age of the universe is estimated to be about 15 billion years).
- It belongs to a class of problems known as *intractable* problems

Conclusion

- Recursion is a very important technique.
- Linear searching v.s. binary searching
- Naïve sort v.s. merge sort
- Good algorithms are better than super computers.
- Computer Science is more than programming!
- The most important computer for any computing professional is between their ears.